

# Closed Forms for Weighted Model Counting

in the Two Variable Fragment

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## Introduction

- Weighted First Order Model Counting  $WFOMC(\Phi, n) = \sum_{\omega \models \Phi} w(\omega)$
- **Probabilistic Inference** in many Probabilistic Logic Frameworks can be cast as WFOMC. But WFOMC is #P!!!.
- Fragments of First Order Logic (FOL) that admit polynomial time WFOMC are called **Domain Lifiable**.
- In this work we revisit the domain liftable fragments of First Order Logic, providing **combinatorial proofs** and **closed for formulas** for WFOMC.

## Universally Quantified Formulas

FOL formula of the form  $\forall x \forall y. \Phi(x, y)$

Example: Simple Undirected Graphs

$$\Phi(x, y) = \neg R(x, x) \wedge (R(x, y) \rightarrow R(y, x))$$

Formula:

$$\sum_{\mathbf{k}, \mathbf{h}} \binom{n}{\mathbf{k}} \prod_{0 \leq i \leq j \leq 2^u - 1} \binom{\mathbf{k}(i, j)}{\mathbf{h}^{ij}} \prod_{0 \leq v \leq 2^b - 1} n_{ijv}^{h_v^{ij}}$$

where  $\mathbf{k} = (k_0, \dots, k_{2^u - 1})$  is a  $2^u$ -tuple of non-negative integers,  $\binom{n}{\mathbf{k}}$  is the multinomial coefficient and for every  $0 \leq i \leq j \leq 2^u - 1$ ,  $\mathbf{h}^{ij} = (h_0^{ij}, \dots, h_{2^b - 1}^{ij})$  is a vector of  $2^b$  integers that sum up to  $\mathbf{k}(i, j)$ .

$$\mathbf{k}(i, j) = \begin{cases} \frac{k_i(k_i - 1)}{2} & \text{if } i = j \\ k_i k_j & \text{otherwise} \end{cases}$$

Key Ideas:

- $k_i$  represents the number of constants that realise the  $i^{\text{th}}$  unary assignment, for example, for a constant  $c$ , either  $R(c, c)$  or  $\neg R(c, c)$  is true. Similarly,  $h_v^{ij}$  represents the number of pairs of constants which realise the  $v^{\text{th}}$  binary assignment, for example  $\{R(c, d), R(d, c)\}$ ,  $\{R(c, d), \neg R(d, c)\}$  etc.
- $\binom{n}{\mathbf{k}} \prod_{0 \leq i \leq j \leq 2^u - 1} \binom{\mathbf{k}(i, j)}{\mathbf{h}^{ij}} \prod_{0 \leq v \leq 2^b - 1} n_{ijv}^{h_v^{ij}}$  represent the possible ways of achieving models with  $k_i$  constants of  $i^{\text{th}}$  unary assignments and  $h_v^{ij}$  constants of  $v^{\text{th}}$  binary assignment.
- $\{n_{ijv}\}$  is a 1 iff the pair of constants with  $i^{\text{th}}$  and  $j^{\text{th}}$  unary assignment and  $v^{\text{th}}$  binary assignment are consistent with  $\Phi$  and 0 otherwise.

## Cardinality Constraints

FOL formula in two variable fragment with constraints on the number of true interpretations of a certain predicate.

Example: Simple Undirected Graphs with only  $l$  edges.

$$\Phi(x, y) = \neg R(x, x) \wedge (R(x, y) \rightarrow R(y, x)) \wedge (|R| = 2l)$$

Formula:

$$\sum_{\mathbf{k}, \mathbf{h} \models \rho} \binom{n}{\mathbf{k}} \prod_{0 \leq i \leq j \leq 2^u - 1} \binom{\mathbf{k}(i, j)}{\mathbf{h}^{ij}} \prod_{0 \leq v \leq 2^b - 1} n_{ijv}^{h_v^{ij}}$$

where  $\rho$  is a constraint on the integer tuples  $\mathbf{k}$  and  $\mathbf{h}^{ij}$ .

Key Ideas:

- $\mathbf{k}$  and  $\mathbf{h}^{ij}$  contain all the cardinality information of each predicate.

## Existential Quantifiers

FOL formula of the form:  $\forall x \forall y. \Phi(x, y) \wedge \bigwedge_{i=1}^q \forall x \exists y. \Psi_i(x, y)$

Example: Simple Undirected Graphs with at least one edge from each node.

$$\forall x y. \neg R(x, x) \wedge (R(x, y) \rightarrow R(y, x)) \wedge \forall x \exists y. R(x, y)$$

Formula:

$$\sum_{\mathbf{k}, \mathbf{h} \models \rho} (-1)^{\sum_i k(P_i)} \binom{n}{\mathbf{k}} \prod_{0 \leq i \leq j \leq 2^u - 1} \binom{\mathbf{k}(i, j)}{\mathbf{h}^{ij}} \prod_{0 \leq v \leq 2^b - 1} n_{ijv}^{h_v^{ij}}$$

Key Ideas:

- Principle of Inclusion Exclusion:  $e_0 = \sum_{l=0}^m (-1)^l s_l$
- $e_0 = \text{FOMC}(\forall x \forall y. \Phi(x, y) \wedge \bigwedge_{i=1}^q \forall x \exists y. \Psi_i(x, y))$
- $s_l = \text{FOMC}(\forall x \forall y. \Phi(x, y) \wedge \bigwedge_{i=1}^q (P_i(x) \rightarrow \neg \Psi_i(x, y)) \wedge (\sum_i |P_i| = l))$

## Counting Quantifiers

Counting Quantifiers are quantifiers of the form  $\exists^{=k}$ ,  $\exists^{\leq}$  and  $\exists^{\geq}$ . Any formula with counting quantifiers can be written as the following model count preserving reduction:

$$\Phi_0 \wedge \bigwedge_{k=1}^q \forall x. (A_k(x) \leftrightarrow \exists^{=m_k} y. R_k(x, y))$$

Example: Simple Undirected  $k$ -Regular Graphs

$$\forall x y. \neg R(x, x) \wedge (R(x, y) \rightarrow R(y, x)) \wedge \forall x \exists^{=k} y. R(x, y)$$

Key Ideas:

- Expressing the counting quantifiers  $\exists^{=k}$  as a set of formulas with existential quantifiers and cardinality constraints.
- Principle of Inclusion Exclusion:  $e_0 = \sum_{l=0}^m (-1)^l s_l$

## Weighted Model Counting

- Weighted First Order Model Counting can be easily obtained by defining a multiplicative real valued weight function on the  $(\mathbf{k}, \mathbf{h})$  vectors.
- Symmetric Weighted Model Counting :

$$w(\mathbf{k}, \mathbf{h}) = \prod_{P \in \mathcal{L}} w(P)^{\mathbf{k}, \mathbf{h}(P)} \cdot \bar{w}(P)^{\mathbf{k}, \mathbf{h}(-P)}$$

- Expressing Count Distributions:

$$w(\mathbf{k}, \mathbf{h}) : (\mathbf{k}, \mathbf{h}) \rightarrow \mathbb{R}$$

## References

1. P. Beame et al., *Symmetric Weighted First-Order Model Counting*, 2015, arXiv: 1412.1505.
2. O. Kuzelka, *Weighted First-Order Model Counting in the Two-Variable Fragment With Counting Quantifiers*, 2020, arXiv: 2007.05619.
3. S. Malhotra et al., *Weighted Model Counting in FO2 with Cardinality Constraints and Counting Quantifiers: A Closed Form Formula*, 2021, arXiv: 2110.05992.